

# Searching the Variability Space to Fix Model Inconsistencies: A Preliminary Assessment

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## ■ ***Software Product Lines (SPL)***

- Families of related systems distinguished by the features
- Extensive success in achieving software reuse
- *Variability*
  - The capacity of software artifacts to change
  - Its effective management is at the core of SPL

## ■ ***Model-Driven Engineering (MDE)***

- Emerging software development paradigm
- Raises the level of abstraction and automates program generation

## ■ Fact

- Increasing research convergence in SPL and MDE that leverages their complementary capabilities

## ■ Challenge

- Maintain consistency between models with variability
  - Checking that certain relations between model elements hold for *all* products of a SPL

## ■ Problem

- Research has focused on inconsistency detection
- Inconsistency fixing has not been fully explored

# Why Search-Based SE?



## ■ Facts

- SPL can employ multiple models simultaneously
  - State charts, sequence diagrams, class diagrams, ...
- Models can have a large number of consistency rules and instances
- SPL usually involve large number of different products

## ■ Because ...

- Large search space
- No unique or optimal fixing solutions

# Our ongoing work ...



- Description
  - Finds fixing locations for single constraint instances
  - Uses basic search algorithm
- Relies on the notion of Safe Composition
  - Programming languages
    - guarantee that *all* programs of a product line are type safe
  - Uses propositional logic

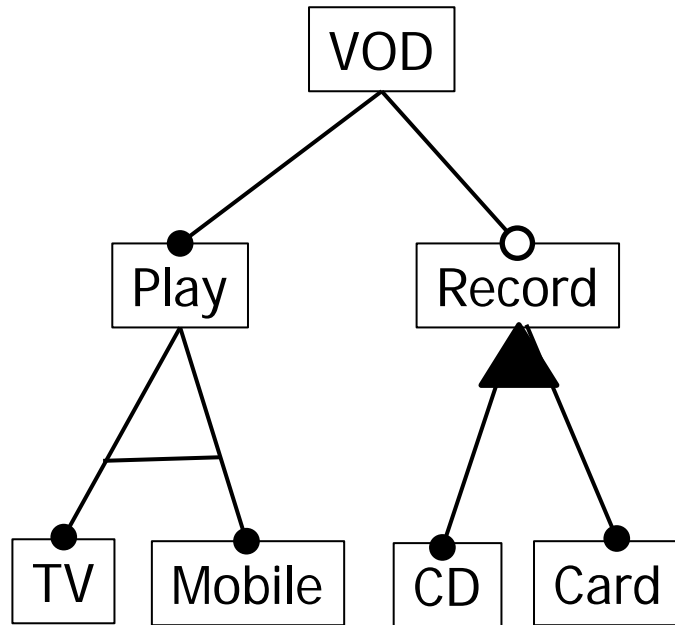
- Intuition
  - Implementation constraint(s) ( $\text{IMP}_f$ ) must follow the product line domain constraints ( $\text{PL}_f$ )
- A SAT solver checks if **one** propositional formula is satisfiable or not
- Our interest is verifying that **all** the product line members satisfy an implementation constraint

$$\neg ( \text{PL}_f \Rightarrow \text{IMP}_f )$$

**Unsatisfiable** = there is no product that violates the constraint

**Satisfiable** = there is at least one product that violates the constraint

# Example $PL_f$



$VOD \Leftrightarrow \text{true} \quad \wedge$

$VOD \Leftrightarrow \text{Play} \quad \wedge$

$\text{Record} \Rightarrow VOD \quad \wedge$

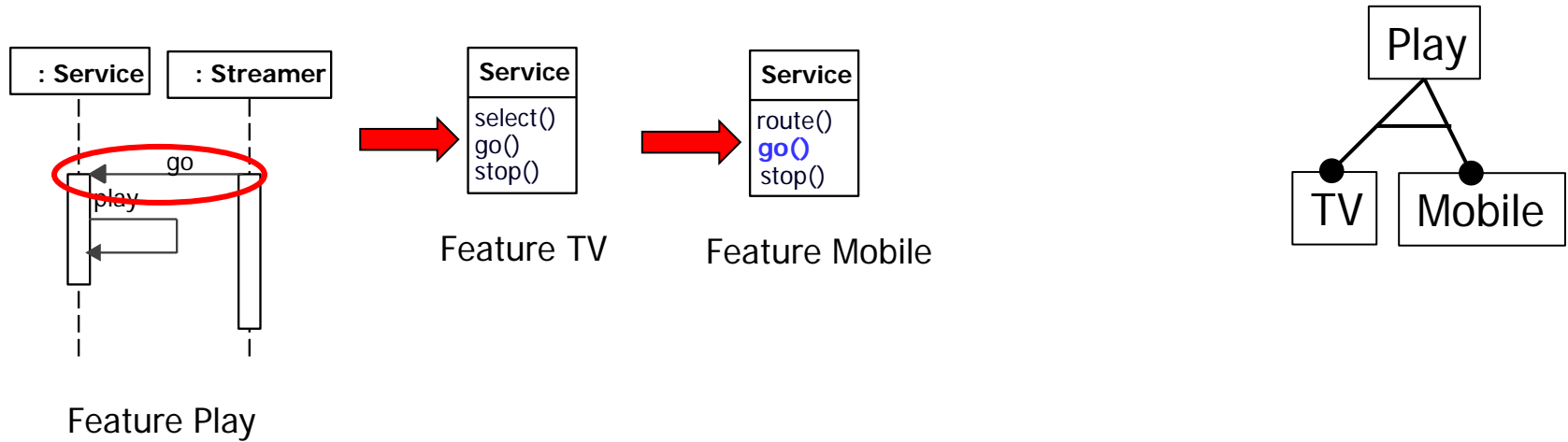
$\text{TV} \Leftrightarrow \neg \text{Mobile} \wedge \text{Play} \quad \wedge$

$\text{Mobile} \Leftrightarrow \neg \text{TV} \wedge \text{Play} \quad \wedge$

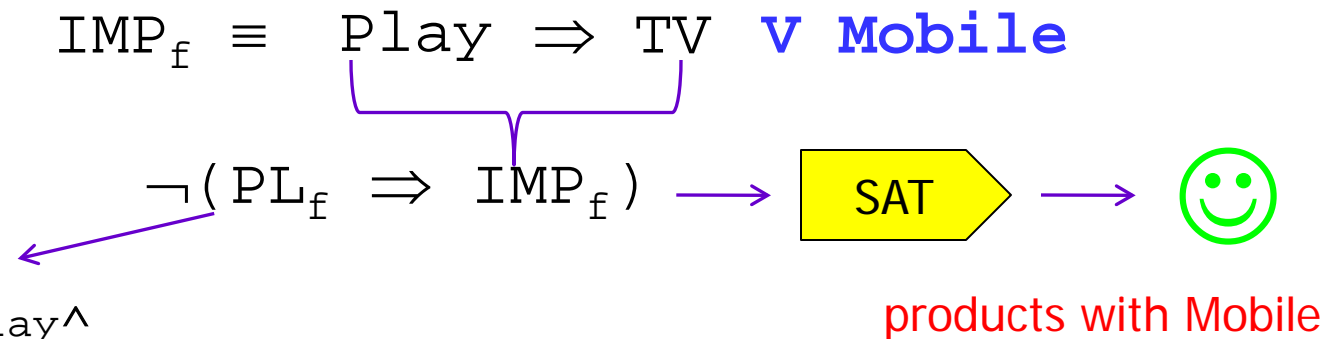
$\text{Record} \Leftrightarrow \text{CD} \vee \text{Card}$

# Requiring Rule Example

- Message action must be defined as an operation in receiver's class



$VOD \Leftrightarrow true \wedge$   
 $VOD \Leftrightarrow Play \wedge$   
 $Record \Rightarrow VOD \wedge$   
 $TV \Leftrightarrow \neg Mobile \wedge Play \wedge$   
 $Mobile \Leftrightarrow \neg TV \wedge Play \wedge$   
 $Record \Leftrightarrow CD \vee Card \wedge$





# Fixing inconsistencies

- **Consistency Rule Instance (CRI)** is a 4-tuple  $[F, RME, TS, FC]$  where:
  - Requiring feature  $F = \text{Play}$
  - Requiring model element  $RME = \text{play}_{\text{msg}}$
  - Set of pairs (feature, required elements)  $TS = \{(TV, \text{play}_{\text{op}})\}$ 
    - Set of features in the pairs of TS as  $TS[\text{feature}]$ .
  - Faulty feature configuration that violates the consistency rule instance.  
 $FC = [\{\text{VOD}, \text{Play}, \text{Mobile}\}, \{\text{TV}, \text{Record}, \text{CD}, \text{Card}\}]$
- **Pair-wise commonality**
  - Operation that receives a feature model P and two features F and G, and returns the number of products that have both features.
- **Question: where to add the required elements?**
  - Intuition: iteratively search fixing configurations choosing first those features with higher commonality → fix more configurations

# Algorithm – Based on BFS

Input: CRI [F,RME,TS, FC] with  $FC \neq \emptyset_{\text{conf}}$ , and  $PL_f$ .

Output: Fixing set FS

$FC' := FC$

$FS := TS[\text{feature}]$

$FSQ.\text{enqueue}(FS)$

while  $FC' \neq \emptyset_{\text{conf}}$  do

$FSQ.\text{dequeue}(FS)$

$G := \text{maxCom}(F, FC', TS, FS)$

$FS := FS \cup G$

$FC' := \text{SafeComposition}(PL_f, F, FS)$

$FSQ.\text{enqueue}(FS)$

end while

return FS

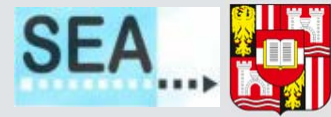
Set of features that guarantees no faulty configurations

Chooses a feature with maximum pair-wise commonality

Applies basic Safe Composition

- Preliminary evaluation
  - Using 60 publicly available feature models, 6-94 features
  - On average fixing sets around 5 elements
- Future work
  - Consider multiple consistency instances
  - Assess other search alternatives
  - Research alternatives for Pair-wise Commonality

# Acknowledgements



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Der Wissenschaftsfonds.

# Questions?